Lecture 2

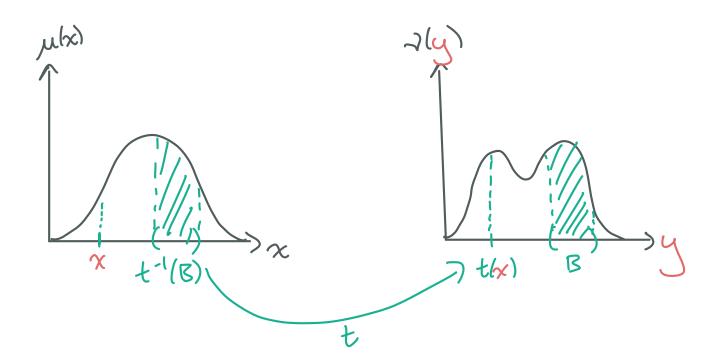
Recall :

Given  $\mu \in P(x)$ ,  $B \in B(x)$ ,  $\mu(B) = ant.$  of dirt in the pile  $\mu$  that lies in B. If  $(x,d) = (\mathbb{R}^d, 1\cdot 1)$  and  $\mu < \lambda$ ,  $d\mu(x) = \mu(x)dx$ and  $\mu(B) = \int_B \mu(x)dx = ant.$  of dirt in B. What does it mean to "rearrange one probability measure to look like another"?  $\forall B \in B(x), t'(B) \in B(x)$ 

Def (transport map): Given:  $\mu^e P(X), \nu \in P(Y),$ a measurable function 5t: X > 4 transports  $\mu$  to  $\nu$  if

 $v(B) = \mu(t^{-1}(B)) \forall B \in B(4)$ 

We call v the pushforward of under t, writing v=t#µ, and we call t the transport map from u to v



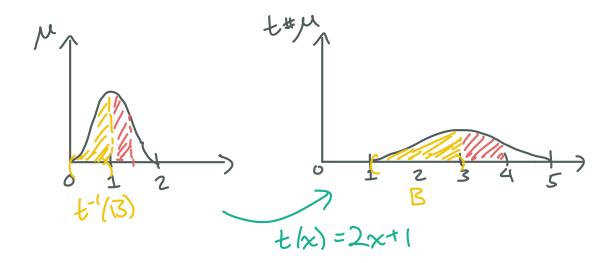
"The amount of mass that & assigns to IS equals the amount of mass sent there from "

Informally, "mass starting at location x in µ is sent to location t(x) in V"

Sanity check: if  $\mu \in P(X)$  and t: X > 4 is measurable, is  $t \neq \mu$  always a prob measure?  $(t \neq \mu)(4) = \mu(t^{-1}(4)) = \mu(X) = 1$ .

Exitranslation/dilation): Suppose  $(\chi, d) = (\mathbb{R}^{d}, 1 \cdot 1)$ . Fix a > 0, be  $\mathbb{R}^{d}$  and  $t(\chi) = a\chi + b$ .

dilation translation Then for any  $\mu \in P(X)$ ,  $t = \mu$  satisfies  $(t * \mu)(B) = \mu(t^{-1}(B)) = \mu(\frac{B-b}{a}) = \mu(\xi = \frac{y-b}{a}; y \in B\xi) + B \in B(A).$ 



Lemma (equiv characterization of transp. map)  
Given 
$$\mu \in P(X)$$
,  $\nu \in P(M \text{ and } t: X \rightarrow Y)$  measurable,  
then  $t * \mu = \nu$  if and only if  
 $S = P(t(X)) d\mu(X) = S = S = P(y) d\nu(y)$  for all  $q: Y \rightarrow \mathbb{R}$   
 $\chi$  measurable,  $q \in L^2(\nu)$   
(\*\*)

Before we prove the lemma, recall:

$$\begin{aligned} & P: First, note that if  $P \text{ is an indicator fn,} \\ & then, using the fact \\ & P(t(x)) = 1_{B}(t(x)) = \begin{cases} 0 & \text{if } t(x) \notin B = 1_{t'(B)}(x), \\ 1 & \text{if } t(x) \notin B \end{cases} \end{aligned}$$$

equation (\*\*) becomes  $S1_{B(1)}d\mu(x) = S1_{E'(B)}(x)d\mu(x) = \mu(E'(B))$ "  $S1_{B(Y)}d\nu(Y) = \nu(B).$ 

Thus: (i) If egn (\*) holds for all 9 measurable with 962<sup>2</sup>(v), then it must hold for

9=1B, so the above remark gives  $\neg (B) = \mu(t^{-1}(B)).$ (ii) If t#µ=v, the above remark shows that (#) holds for all indicator functions. Thus, Et) for all q meas w/ lel<sup>2</sup>(v) implies t # u=v.

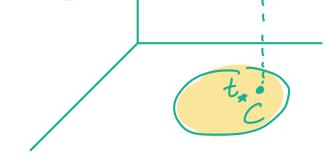
Now, assume 2# n=v. We have (#) for all indicator fns, 9=18. Furthermore, by linearity of the integral, (#) holds for all simple functions 9. nonneg integrable Next, suppose quis a bold, mass fr. Choose a sequence qui of simple fris so that qui q pointwise. Thus by the dominated convergence theorem, S P(t(x)) duk = "" S Pult(x)) dulx)  $= \lim_{n \to \infty} \int \Pr(y) d v(y)$ = SP(y)dv(y)

Thus, 6th holds for all 9 bdd, meas. nonneg integrable Noxt, suppose q is a nonnegative, meas fin in 2<sup>1</sup>(v), and define  $f(x) = Q(x) \wedge n = min(Q(x), n)$ Then, by the monotone convergence theorem, S P(t(x)) detet in S Putt(x)) dulx) Elim Jen (y)dvly = SP(y)dvly)Finally, for q an arbitrary meas fn in L<sup>2</sup>(v), the result holds by writing,  $\varphi(x) = \varphi_{+}(x) - \varphi_{-}(x) = \varphi(x) \vee O - (-\varphi) \vee O \square$ max(Q(x), 0)Now, we know what it means to rearrange one measure to look like another", or, more precisely, to transport one measure to another. Back to original question: how can this be done in the most efficient way?

Monae's Optimal Transport Problem; Giver M, V E P(X), Solve how for dift is moved Jdbx, t(x) dp(x) min t: X->X measurable how much dirt 七#ル=ン effort to rearrange in to look like ~ via the transport map t

Throughout the course, we'll see many optimization problems of this form: min F(t) teC F. constraint set

Mental image: +(++)



Unfortunately, Monge's problem is a horrible optimization problem!

Sudakov 1979, Ambrosio and Pratelli 2001 Evans and Gangbo 1999

Reasons the Monge Problem is difficult: Difficulty #1: the constraint set can be empty. That is, given  $\mu, \nu \in P(x)$ , there doesn't necessarily exist any time as s.t.  $t*\mu=\nu$ . Recall: S<sub>x</sub> is the probability measure S<sub>x</sub>(B)={0 if x #B (1 if x \*B



 $\frac{1}{2} = \frac{1}{2} \sqrt{2} \sqrt{2} \sqrt{2}$  If  $t^{\#}\mu = \sqrt{2}$ , then  $\lambda(B \cap [0,0]) = \sqrt{2}(B) = \mu(t^{-1}(B)) = \begin{cases} 0 & \text{if } t(\frac{1}{2}) \notin B \\ 1 & \text{if } t(\frac{1}{2}) \in B \end{cases}$ 

There is no such I for which this holds.

Heuristically, the problem is that a transport map t sends all mass starting at a location x. to t(x.). In particular, mass cannot split.

On the other hand, note that  $t(x)=\frac{1}{2}$ satisfies

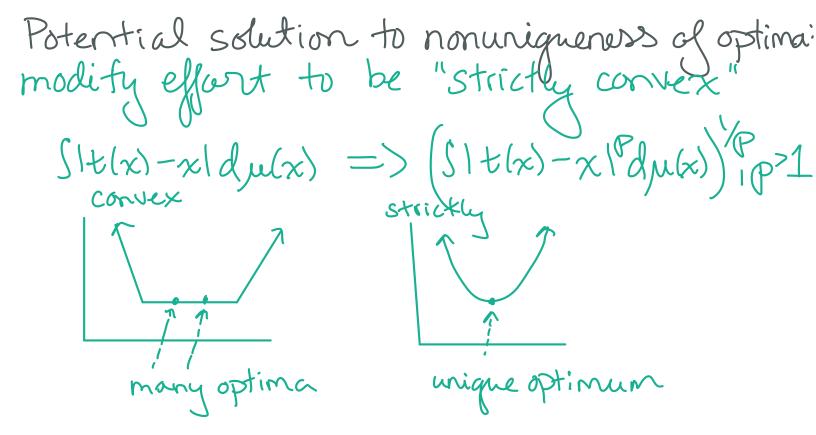
 $(t + v)(B) = -i(t'(B)) = S_v(B) \quad if \neq B = S_1(B) = u(B).$  $v(\chi) \quad if \neq eB$ 

Iwo potential solutions to empty constraint set: (a) don't allow source measure to concentrate mass on "small sets" (b) instead of considering transport maps, consider transport plans.

Difficulty #2: Solutions may not be unique. That is, given unv EP(x), there may exist multiple, distinct optimal transport maps. Ex: "books on shelf"  $1 + \frac{1}{1} +$ 

Consider  $t_0(x) = x + \frac{1}{4}$  "shifting all bookstoright"  $t_2(x) = (x + 1) \quad \text{if } x \in [0, \frac{1}{4})$  "shift first  $(x) \quad \text{otherwise}$  book to end"

Exercise: to# u=v and ti# u=v, so both to and to belong to the constraint set. Fact (will show later): to and t, are both optimal transport maps.



Difficulty #3: The constraint set is nonconvex

Recall:

 $\mathcal{D}_{ef}$ : A subset C of a vector space X is <u>convex</u> if,  $\forall x_{0,X_{1}} \in C$ ,

 $\chi_{\alpha} := (1 - \alpha)\chi_{0} + \alpha\chi_{1} \in C, \quad \forall \alpha \in [0, 1].$ 

Generally, in optimization, we want our constraint set 'C to be convex, since our normal strategy is to take an initial quess, perturbilit, and see if the objective function decreases.

